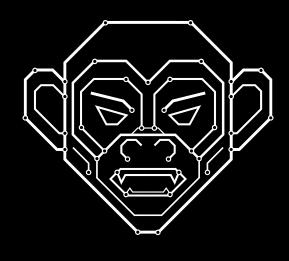
AFNOM Session 5

Binary Exploitation/PWN



A - F - N - O - M

Society Recap

- We learn hacking skills and do CTFs
- No membership just turn up!
- Come to the pub after sessions + occasional events (films / games)
- Ask questions!! We're all here to learn from each other :-)

AFNOM Values

- Inclusive everyone is welcome, regardless of race, ethnicity, gender, gender identity, sexual orientation, age, class, physical ability, nationality, and text-editor preference.
- **Supportive** we're not here to compete with each other; we've all got things to learn and teach.
- Ambitious we all want to advance our cybersecurity skills.
- Respectful let's be kind and humble.

Preface

- Some things we cover could be interpreted as or used immorally or illegally
- We are here to learn cool stuff and have fun!

"Don't do crimes" - Dr lan Batten

- Good lecturer
- Sound advice
- We are the **ethical** hacking society

"Don't be evil" - Google

- Somewhat ironic
- They stopped using it in 2015
- But we are not multi conglomerate mega empire so we can reserve the right to morals

Upcoming schedule

- Today Binary Exploitation/PWN
- 12th November CSAW Recap + Hack n Chill
- 19th November Network Forensics
- **26th November** Cryptography
- **Future** ... Any suggestions? Let us know!

What is PWN?

- Exploiting the functionality of a binary (native program) to circumvent protections, gain access/permissions, etc.
- Redirecting control flow, running your own code, or dropping into a shell
- One of the things you might imagine when you think of 'hacking'

Expectation



"011101010110101001.....Im in"





Today's Session

- Intro to PWN
- Deep and technical category
- Want to get you started with enough to give it a go!

Covered today

- Assembly & the stack how code runs at a low level (on the x86 architecture)
- Smashing the stack, taking control, and a little more

Not today

- Other architectures (ARM, RISCV, etc)
- How to do harder stack attacks (ROP)
- The heap + heap attacks 🤮
- Maybe in a future session though!

Kickstart - the basics

Assembly

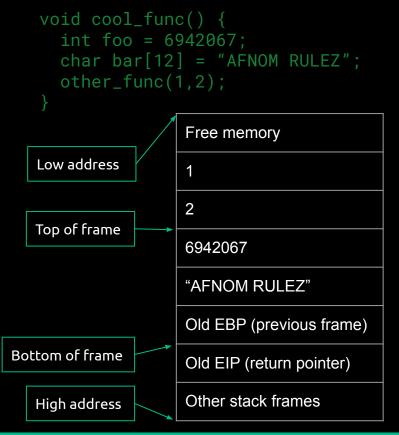
- Computers can't just run code raw
- It needs to be processed and converted (compiled) into machine code instructions that the computer knows
- Machine code can then be interpreted and represented as assembly code
- Changes between architectures
- E.g. 'mov ebx, eax', 'add ebp, 0x10'

The stack

- Section of memory used for storing local variables, working data, etc.
- Each function gets its own 'frame'
- Frame size calculated at compile time
- Grows down memory starts at higher memory address - frames are stacked at decreasing addresses
- However, frames filled 'up' values written in increasing memory addresses like you might expect
- Sometimes stack stuff is shown low->high, sometimes high->low
- Important to understand and identify this to avoid confusion!

Stack Layout - Variables, parameters and control flow

- When a function is called, a new frame is created
- Return address (next instruction in caller function) and previous frame pointer are at the bottom of the frame
- Local variables/data written to frame from top to bottom
- Calling conventions
 - 32-bit: Arguments are pushed to the stack
 - 64-bit: Arguments put in registers:
 RDI, RSI, RDX, RCX, R8, R9, and then stack



Diving in - Common attacks

- Buffer overflow
 - Writing more data than intended to a memory location
- Ret2Win
 - Overwrite return pointer to jump to a specific function
- Ret2LibC
 - Ret2Win but specifically jumping into the C standard library
- Leaking stack memory
 - Finding memory addresses (stack base, code base, libc base, etc.)
- Shellcode
 - Inserting and executing your own code/assembly in your inputs
- Return Oriented Programming (ROP)
 - Reusing program code to build/execute your own

Protections (checksec)

- RELRO RELocation Read-Only
 - Library function pointers cannot be overwritten
- Canary Stack Canaries
 - A static value is placed before the stack frame to detect if the frame has been corrupted
 - Prevent buffer overflows
- NX / DEP Not eXecutable / Data Execution Prevention
 - Memory pages cannot be loaded as both writable and executable
 - Prevent shellcode
- PIE Position Independent Execution
 - The address of the binary base is randomized
 - Prevent ret2win/rop
- ASLR Address Space Layout Randomisation
 - Randomizes the location of stack, heap, shared objects, etc.
 - Prevent ret2libc/rop

Smashing the stack - Buffer Overflow

- The size of values on the stack are fixed and pre-calculated based on the code (type, defined length, value)
- Inputting/moving data should use these bounds but aren't enforced
- Using functions like 'gets' or even 'fgets' with the wrong size can let you write past the end of the buffer
- Can overwrite other variables or even return pointer
- Basis for most PWN attacks

```
void get_input()
  char buf[15];
  gets(buf);
```

Normal input: "Hello World!"

<u>"Hello World!"</u>
Free memory
"Hello World!"
Old EBP (previous frame)
Old EIP (return pointer)
Other stack frames

Free memory
"аааааааааааааа"
"aaaa"
"aaaa"
Other stack frames

Return pointer overwritten *program crashes*

Taking control - Ret2Win/Ret2LibC

- Use buffer overflows for more than just crashing the program
- Instead of just overflowing with random characters, why not give valid return address
- Can change what code runs after the current function!
- x86 is little endian, so need to be input in reverse order
- Not all hex values are printable, use a command like whilehex to input (shown later)

```
void win() {
  printf("win!\n");
}
void main() {
  char buf[10];
  gets(buf);
}
```

Compile 64-bit with no canary + no PIE

win() is at 0x401156

Let's attack this program!

Step 1: overflow buffer + old frame pointer Step 2: reverse + overwrite return pointer Step 3: profit??

Leaking stack values

- PIE randomises the memory base of the program
- ASLR randomises the base of the stack
- Canaries are random values
- On a remote binary you can't attach a debugger
- The printf vulnerability can be used to leak many stack values (canarys, addresses, variables)
- Given there is enough space for a long enough format string

- Common leaking method uses 'printf' calls that print user input
- Correct use: printf(formatstring, userinput)
- Vulnerable use: printf(userinput)
- Will treat your input as the format string
- Input string that would format arguments e.g. "%p%p%p%p%p"
- Then printf will start taking values for arguments
- 32-bit: read stack values
- **64-bit**: registers then stack values

Complete freedom - Shellcode & ROP

Shellcode

- No function doing what you want to do already?
- Insert your own assembly code in your input
- Then set return pointer to the current stack frame to run it!
- NX/DEP generally prevents this as makes stack non-executable
- But there are definitely ways around this!

Return Oriented Programming (ROP)

- Similar to shellcode, build your own assembly code
- Build out of 'gadgets' small blocks of instructions ending in return instructions
- Tools like Ropper can auto find and filter gadgets, checking at every byte offset
- Given a large enough binary can build a ROP chain for anything!

Tooling

- Ghidra/IDA/Binary Ninja
 - Disassemblers/Decompilers for viewing and analysing binaries
- GDB
 - Debugging binaries
 - Good for testing your payloads and looking at memory
 - Recommend PwnDBG/GEF
- Python pwntools
 - Library for automating solves
 - Saves you manually copying and formatting payloads
- Our lord and saviour, whilehex:

```
while read -r line; do echo -e $line; done | ./[binary]
```

- Converts hex bytes written as \xXX into the actual byte values, input non printable characters
- Would recommending aliasing in your terminal so can easily pipe into your binary or gdb

Right Now

sessions.afnom.net

- Hang around at 5pm for pub social!
- Come back next week for CSAW Recap + Hack n Chill (PWN)!